

Kurs: MAN321/TMV026 Ändliga automater och formella språk

Plats: M-huset

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No help documents

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*The questions can be answered in english or in swedish.*

total 30;  $\geq 13$ : 3,  $\geq 19$ : 4,  $\geq 25$ : 5

total 30;  $\geq 13$ : G,  $\geq 21$ : VG

1. Let  $\Sigma$  be an alphabet. What is, mathematically, a deterministic finite automaton on  $\Sigma$  (1p)? Explain what is the language determined by such a finite automaton (1p). Explain why such a language is a context-free language (1p).
2. Minimize the following automaton (2p)

	a	b
$\rightarrow 0$	1	3
1	0	3
2	1	4
*3	5	5
4	3	3
*5	5	5

3. Build a NFA with 3 states that accepts the language  $\{ab, abc\}^*$ . (2p)
4. Build a DFA corresponding to the regular expression  $(ab)^* + a^*$ . (3p)
5. Let  $\Sigma$  be  $\{0, 1\}$ . We recall that the regular expressions are on  $\Sigma$  are given by the grammar

$$E ::= 0 \mid 1 \mid \emptyset \mid \epsilon \mid E + E \mid EE \mid E^*$$

Give a regular expression  $E$  such that

$$L(E) = \Sigma^* - L(10(0+1)^*) \quad (2p)$$

6. Build a DFA that recognizes exactly the word in  $\{0, 1\}^*$  ending with the string 1110. (2p)
7. Is the following grammar ambiguous? Why (2p)?

$$S \rightarrow AB \mid aaB, \quad A \rightarrow a \mid Aa, \quad B \rightarrow b$$

Construct an unambiguous grammar which is equivalent to this grammar (2p).

8. Consider the grammar

$$S \rightarrow aaB, \quad A \rightarrow bBb \mid \epsilon, \quad B \rightarrow Aa$$

Show that the string *aabbabba* is *not* in the language generated by this grammar (3p).

9. Find context-free grammars for the languages

(a)  $L = \{ a^n b^n c^k \mid n \leq k \}$  (1p)

(b)  $L = \{ a^n b^m \mid n \neq m \}$  (1p)

10. Let  $L, M, N$  be languages on an alphabet  $\Sigma^*$  (that is, we have  $L, M, N$  subsets of  $\Sigma^*$ ). Explain why we have  $L(M \cap N) \subseteq LM \cap LN$  (2p). Give an example showing that we do not have  $LM \cap LN \subseteq L(M \cap N)$  in general (2p).
11. Let  $\Sigma$  be an alphabet and let  $L_1, L_2$  be subsets of  $\Sigma^*$ . Assume that  $L_1 \cap L_2 = \emptyset$  and  $L_1$  is finite and that  $L_1 \cup L_2$  is regular. Can we deduce that  $L_2$  is regular (3p)?